# Review: Data Representation and Boolean operators in C 

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## Binary Representation: Summary

- Every storage locations stores a finite sequence of bits
- 8-bit, 16-bit, 32-bit etc.
- The same bit string can mean different things depending on how the program wants to look at it.

| Address | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 35 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |
| 36 | 1 | 0 | 0 | 0 | 0 | 1 | 1 | 1 |$\quad$ 2C: -127

## Unsigned \& Signed Numeric Values

| $X$ | $\mathrm{~B} 2 \mathrm{U}(X)$ | $\mathrm{B} 2 \mathrm{~T}(X)$ |
| :---: | :---: | :---: |
| 0000 | 0 | 0 |
| 0001 | 1 | 1 |
| 0010 | 2 | 2 |
| 0011 | 3 | 3 |
| 0100 | 4 | 4 |
| 0101 | 5 | 5 |
| 0110 | 6 | 6 |
| 0111 | 7 | 7 |
| 1000 | 8 | -8 |
| 1001 | 9 | -7 |
| 1010 | 10 | -6 |
| 1011 | 11 | -5 |
| 1100 | 12 | -4 |
| 1101 | 13 | -3 |
| 1110 | 14 | -2 |
| 1111 | 15 | -1 |

- Equivalence
- Same encodings for nonnegative values
- Uniqueness
- Every bit pattern represents unique integer value
- Each representable integer has unique bit encoding
- $\Rightarrow$ Can Invert Mappings
- $\mathrm{U} 2 \mathrm{~B}(x)=\mathrm{B} 2 \mathrm{U}^{-1}(x)$
- Bit pattern for unsigned integer
- $\mathrm{T} 2 \mathrm{~B}(x)=\mathrm{B}^{-1}(x)$
- Bit pattern for two's comp integer


## Relation between Signed \& Unsigned

Two's Complement


$$
u x= \begin{cases}x & x \geq 0 \\ x+2^{w} & x<0\end{cases}
$$

## Basic Logic Operations/Gates

- AND:
- Equivalent notations: A AND $\mathrm{B}=\mathrm{A} \cdot \mathrm{B}=\mathrm{A} \wedge \mathrm{B}$
- OR
- Equivalent notations: $A$ or $B=A+B=A \vee B N O T$

- NOT
- Equivalent notations: $\operatorname{not} A=A^{\prime}=\bar{A}$

- XOR
- Equivalent Notations: A XOR B = A ^ B
- Other common logic operations:
- NAND = NOT AND $\square$
- NOR = NOT OR



## Next...a little bit of "reality"

- look at how some of the concepts we have studied take shape in 'real life'
- C programming language


## Data Representations

- Sizes of C Objects (in Bytes)
- C Data Type Compaq Alpha
- int

4
long int
8

- char
- short 1
- float
- double
- long double
- char *
- Or any other pointer

Typical Intel IA32/IA64
4
8
8
1
2
2
4
8
:
4/8 (64 bit needs 8$)$

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## Signed vs Unsigned in C

- C allows int to be defined as unsigned or signed (!!)
- Constants
- By default are considered to be signed integers
- Unsigned if have "U" as suffix

OU, 4294967259 U

- Casting - nasty stuff!!! Or is it fun ??
- Explicit casting between signed \& unsigned same as U2T and T2U
int tx, ty;
unsigned ux, uy;
tx $=$ (int) $u x$;
$u y=$ (unsigned) ty;
- Implicit casting also occurs via assignments and procedure calls
tx = ux;
uy = ty;


## Casting Signed to Unsigned

- C Allows Conversions from Signed to Unsigned

```
short int x = 15213;
unsigned short int ux = (unsigned short) x;
short int y = -15213;
unsigned short int uy = (unsigned short) y;
```

- Resulting Value
- No change in bit representation
- Nonnegative values unchanged
- $u x=15213$
- Negative values change into (large) positive values
- uy $=50323$
- Casting Surprises in expression evaluation
- If you mix signed and unsigned then signed cast to unsigned....and unexpected results in comparisons (>, < etc.)

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## Why Should I Use Unsigned?

- Don't Use Just Because Number Nonzero
- Easy to make mistakes
for (i $=$ cnt-2; $i>=0$; $i--$ )
$a[i]+=a[i+1]$;
- Do Use When Performing Modular Arithmetic
- Multiprecision arithmetic
- Other esoteric stuff
- Do Use When Need Extra Bit's Worth of Range
- Working right up to limit of word size


## Logical Operations in C

- C supports both bitwise and boolean logic operations
- x\&y bitwise logic operation
- $x \& \& y$ boolean operation: output is boolean value
- What's going on here?
- In boolean operation the result has to be TRUE (1) or FALSE (0)
- Treats any non-zero argument as TRUE and returns only TRUE (1) or FALSE (0)
- In C: logical operators do not evaluate their second argument if result can be obtained from first
- a \&\& $5 / \mathrm{a}$ can we get divide by zero error?


## Bitwise Logical Operators

- View $n$-bit number as a collection of $n$ logical values
- operation applied to each bit independently
- Number operated on is an n-bit number
- Operation being performed is logical operation on each bit
- Masking operations
- If we are only interested in last 8 bits of a 32 bit number $X$, how to extract this?
- X \& 0xFF (0xFF is notation for hex number FF)
- Zero out the most significant 24 bits; value of least significant 8 bits is same as the value of these in $X$
- xABCD27A4 \& 0xFF = xA4 (in 32 bits: 0x000000A4)
- $X \& 0 \times 1=0$ if $X$ is even and $=1$ if $X$ is odd


## Bitwise Logical Operations

- View n-bit field as a collection of $n$ logical values
- Apply operation to each bit independently
- Bitwise AND: useful for clearing bits 11000101
- AND with zero $=0$
- AND with one = no change

AND $\quad \frac{00001111}{00000101}$

- Bitwise OR: useful for setting bits
- OR with zero = no change
- OR with one = 1

11000101

- Computers don't support individual bits as a data type
- Just use least significant of $n$-bit integer
- Integers are generally more useful
OR $\frac{00001111}{11001111}$


## Bitwise Operators in C

- Can only be applied to integral operands
- that is, char, short, int and long
- (signed or unsigned)
\& Bitwise AND
| Bitwise OR
^ Bitwise XOR
<< Shift Left
>> Shift Right
~ 1's Complement (Inversion)


## Bitwise AND

- Bitwise AND: 0101 AND 0110 in C: (5 \& 6)
- 0100
- Bitwise OR: 0101 OR 0110 in C:(5 | 6)
- 0111
- Bitwise NOT: NOT 0101 in C: ~5
- 1010
- Bitwise XOR: 0101 XOR 0110 in C: 5^6
- 0011
- Bitwise NAND - no C operator, therefore
- 0101 NAND 0110 in C: ~ (5 \& 6)
- Bitwise NOR - no C operator, therefore
- 0101 NOR 0110 in C: ~(5 | 6)


## Shift Operations

- $\mathrm{x}=01100001$ and $\mathrm{y}=2$ (using 8-bit numbers)
- $z=10100001$
- x >> y
- x right shifted y bit positions, sign extended/arithmetic shift - Sign bit shifted into positions vacated by shifted bits
- $x=01100001 \quad y=2$ (using 8-bit numbers)
- x >> y = 00011000
- z >> y =11101000
- $x \ll y$
- x left shifted y bit positions, zero placed in positions vacated by shifted bits
- $x \ll y=10000100$
- $z \ll y=10000100$
- In C, $x, y$ are 32 bit numbers:
- What is $F=(x \gg 31) \& 0 x 1$


## Boolean Relational Operators

- What is the semantics of:
- If $(x==0)$ then ......
- how many outcomes for $(x==0)$ ?


## - Concept of boolean operators

- Apply logic operators, but treat input and output as boolean variables - Only 1 or 0 (True or False) values for entire variable
- But input strings can be n-bits long?
- Treat entire string as ONE boolean variable - How?

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## Logical Operations in C

- ! Logical NOT
-! $x$
$0!x=0$ if $x$ is non-zero, ! $x=1$ if value of $x$ is zero
- \&\& Logical AND
- $x \& \& y$
$\circ x \& \& y=1$ if value of $x$ is not zero and value of $y$ is not zero
$\circ x \& \& y=0$ if both $x$ and $y$ are zero
- || logical OR
- $x \| y$
$\circ x \| y=1$ if at least one of $x, y$ are not zero
o $x \| y=0$ if both $x, y$ are zero


## Examples

- 8 bit numbers, $f=7, g=8$
- $f=00000111 \quad g=00001000$
- $h=(f \& g)(b i t w i s e ~ A N D) . .$. .
-h= 00000000
- $h=(f \& \& g)(l o g i c a l$ AND)...
- $\mathrm{h}=1$
- $!h=0$ since $h$ is non-zero
- $h=(f \mid g) \quad$ (bitwise $O R) \ldots h=$ ?
- $h=(f| | g)(\operatorname{logical} O R) \ldots h=$ ?
- $h=(\sim f \mid \sim g) \ldots h=$ ?
- $h=(!f \& \&!g) \ldots h=$ ?


## Byte-Oriented Memory Organization

- Programs Refer to Virtual Addresses
- Conceptually very large array of bytes
- Actually implemented with hierarchy of different memory types
- SRAM, DRAM, disk
- Only allocate for regions actually used by program
- In Unix and Windows, address space private to particular "process"
- Program being executed
- Program can clobber its own data, but not that of others
- Compiler + Run-Time System Control Allocation
- Where different program objects should be stored
- Multiple mechanisms: static, stack, and heap
- In any case, all allocation within single virtual address space


## Encoding Byte Values

- Byte = 8 bits
- Binary $00000000_{2}$ to $11111111_{2}$
- Decimal: $0_{10}$ to $255_{10}$
- Hexadecimal $00_{16}$ to $\mathrm{FF}_{16}$
- Base 16 number representation
- Use characters ' 0 ' to ' 9 ' and ' A ' to ' F '
- Write FA1D37B ${ }_{16}$ in C as 0xFA1D37B
- Or 0xfald37b

| $4^{e^{t}} p^{e^{c i n}} \operatorname{Bn}^{n^{2}} n^{2}$ |  |  |
| :---: | :---: | :---: |
| 0 | 0 | 0000 |
| 1 | 1 | 0001 |
| 2 | 2 | 0010 |
| 3 | 3 | 0011 |
| 4 | 4 | 0100 |
| 5 | 5 | 0101 |
| 6 | 6 | 0110 |
| 7 | 7 | 0111 |
| 8 | 8 | 1000 |
| 9 | 9 | 1001 |
| A | 10 | 1010 |
| B | 11 | 1011 |
| C | 12 | 1100 |
| D | 13 | 1101 |
| E | 14 | 1110 |
| F | 15 | 1111 |

## Machine Words

- Machine Has "Word Size"
- Nominal size of integer-valued data
- Including addresses
- Some current machines are 32 bits (4 bytes)
- Limits addresses to 4GB
- Becoming too small for memory-intensive applications
- Most (and all Higher-end) systems are 64 bits ( 8 bytes)
- Potentially address $\approx 1.8 \times 10^{19}$ bytes
- Machines support multiple data formats
- Fractions or multiples of word size
- Always integral number of bytes


## Word-Oriented Memory Organization

- Addresses Specify Byte Locations
- Address of first byte in word
- Addresses of successive words differ by 4 (32-bit) or 8 (64-bit)

| 32-bit <br> Words | 64-bit <br> Words | Bytes | Addr. |
| :---: | :---: | :---: | :---: |
|  |  |  | 0000 |
| Addr $=$ |  |  | 0001 |
| 0000 |  |  | 0002 |
|  | Addr $=$ |  | 0003 |
|  | 0000 |  | 0004 |
| Addr $=$ |  |  | 0005 |
| 0004 |  |  | 0006 |
|  |  |  | 0007 |
|  |  |  | 0008 |
| Addr $=$ |  |  | 0009 |
| 0008 |  |  | 0010 |
|  | Addr $=$ $=$ |  | 0011 |
|  | 0008 |  | 0012 |
| Addr |  |  | 0013 |
| 0012 |  |  | 0014 |
|  |  |  | 0015 |

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## Byte Ordering

- How should bytes within multi-byte word be ordered in memory?
- Conventions
- PowerPC (old Mac' s) are "Big Endian" machines
- Least significant byte has highest address
- Big end first
- Intel x86, PC's are "Little Endian" machines
- Least significant byte has lowest address
- Little end first
- Most network protocols use Big Endian
- The terms big-endian and little-endian come from Jonathan Swift's eighteenthcentury satire Gulliver's Travels. The subjects of the empire of Blefuscu were divided into two factions: those who ate eggs starting from the big end and those who ate eggs starting from the little end.


## Byte Ordering Example

- Big Endian
- Least significant byte has highest address
- Little Endian
- Least significant byte has lowest address
- Example
- Variable x has 4-byte representation $0 \times 01234567$
- Address given by $\& x$ is $0 \times 100$



## Representing Integers

- int $A=15213$;
- int $B=-15213$;
- long int $C=15213$;

| Decimal: | 15213 |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| Binary: | 0011 | 1011 | 0110 | 1101 |  |
| Hex: 0000 | 3 | B | 6 | D |  |
| Decimal: | -15213 |  |  |  |  |
| Hex: | FFFF | C | 4 | 9 | 3 |

-Little endian layout for A :
-For B
-For C

- Big endian layout for $A$ :
-For B:
-For C:


## Representing Integers

- int $A=15213$;
- int $B=-15213$;
- long int $C=15213$;


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## Why this discussion of Bit manipulation operations in C.....Project 2!

- Project 2: Given a set of functions, each of which does not use conditional statements and implements some bit manipulation function, determine the function being implemented.
- Rewrite the code to provide an equivalent more readable code using any $C$ operators including conditional statements.
- Why is this useful - some functions can be executed much quicker if they can re-written using bit manipulation operations

